

1 Conceptual Check

Order the following big-O runtimes from most to least efficient:

$O(n \log n)$, $O(1)$, $O(2^n)$, $O(n^2)$, $O(\log n)$, $O(n)$, $O(n!)$

_____ \subset _____ \subset _____ \subset _____ \subset _____ \subset _____

Are the statements in the right column true or false? If false, correct the asymptotic notation (Ω , Θ , O). $\Omega(\cdot)$ is the opposite of $O(\cdot)$, i.e. $f(n) = \Omega(g(n)) \iff g(n) = O(f(n))$.

$f(n) = 100$	$g(n) = 1$	$f(n) \in \Omega(g(n))$	_____
$f(n) = n^2 + n$	$g(n) = 0.1n^2$	$f(n) \in \Theta(g(n))$	_____
$f(n) = 2^n$	$g(n) = 2^{2n} + 100$	$f(n) \in \Theta(g(n))$	_____
$f(n) = n^{100}$	$g(n) = 2^n + n \log n$	$f(n) \in O(g(n))$	_____
$f(n) = 3 \log n$	$g(n) = n^2 + n + \log n$	$f(n) \in \Omega(g(n))$	_____
$f(n) = n \log n$	$g(n) = (\log n)^2$	$f(n) \in O(g(n))$	_____

2 Analyzing Runtime

Give the worst case runtime in $\Theta(\cdot)$ notation. Extra: Give the best case runtime in $\Theta(\cdot)$.

A. Use M and N in your result. `bump()` is a constant time function that returns a boolean.

```

1 int j = 0;
2 for (int i = 0; i < N; i += 1) {
3     for (; j < M; j += 1) {
4         if (bump(i, j))
5             break; // terminates only the inner loop
6     }
7 }
```

B. Use N in your result.

```

1 public static boolean mystery(int[] arr) {
2     arr = bilbosort(arr); // creates sorted copy of arr in  $\Theta(N \log N)$  time
3     int N = arr.length;
4     for (int i = 0; i < N; i += 1) {
5         boolean x = false;
6         for (int j = 0; j < N; j += 1) {
7             if (i != j && arr[i] == arr[j])
8                 x = true;
9         }
10        if (!x)
11            return false;
12    }
13    return true; } // } on same line for vertical space reasons
```

C. Use N in your result, where N is the length of `arr`. Assume `arr` is a sorted array of unique elements. Say we call `mystery2(arr, 0, arr.length)`.

```
1 public static int mystery2(int[] arr, int low, int high) {
2     if (high <= low)
3         return -1;
4     int mid = (low + high) / 2; // (later, see http://goo.gl/gQI0FN )
5     if (arr[mid] == mid)
6         return mid;
7     else if (mid > arr[mid])
8         return mystery2(arr, mid + 1, high);
9     else
10        return mystery2(arr, low, mid);
11 }
```

Amanda's Additional for Awesome Asymptotic And Algorithmic Allstars with Alliteration: What are `mystery()` and `mystery2()` doing? Assuming an `int` can appear in `arr` at most twice, can you rewrite `mystery()` with a better runtime?

3 Optimizing Algorithms (Extra Problem)

Given an integer x and a **sorted** array `A[]` of N distinct integers, design an algorithm to find if there exists indices i and j such that $A[i] + A[j] == x$.

Let's start with the naive solution.

```
1 public static boolean findSum(int[] A, int x) {
2     for (int i = 0; i < _____; i++) {
3         for (int j = 0; j < _____; j++) {
4             if (_____) {
5                 _____;
6             }
7         }
8     }
9     return false;
10 }
```

Can we do this faster? Hint: Does order matter here?

```
public static boolean findSumFaster(int[] A, int x){
```

```
}
```

What is the runtime of both these algorithms?